

**AMENDMENTS TO THE CLAIMS**

Pursuant to 37 C.F.R. § 1.121 the following listing of claims will replace all prior versions, and listings, of claims in the application.

**Listing of the Claims**

1 – 2. (Canceled)

3. (Withdrawn) A method for establishing a common key for a group of at least three subscribers, the method comprising:  
 generating by each subscriber  $T_i$  of the at least three subscribers a respective message  $N_i = (g^{z_i} \bmod p)$  from a publicly known element  $g$  of large order of a publicly known mathematical group  $G$  and a respective random number  $z_i$  and sending the respective message from the respective subscriber to all other subscribers  $T_j$  of the at least three subscribers, each respective random number  $z_i$  being selected or generated by the respective subscriber  $T_i$ ;  
 generating by each subscriber  $T_i$  a transmission key  $k^{ij}$  from the messages  $N_j$  received from the other subscribers  $T_j$ ,  $j \neq i$ , and the respective random number  $z_i$  according to  $k^{ij} := N_j^{z_i} = (g^{z_j})^{z_i}$ ;  
 sending by each subscriber  $T_i$  the respective random number  $z_i$  in encrypted form to all other subscribers  $T_j$  by generating the message  $M_{ij}$  according to  $M_{ij} := E(k^{ij}, z_i)$ ,  $E(k^{ij}, z_i)$  being a symmetrical encryption algorithm in which the data record  $z_i$  is encrypted with the transmission key  $k^{ij}$ ; and  
 determining a common key  $k$  by each subscriber  $T_i$  using the respective random number  $z_i$  and the random numbers  $z_j$ ,  $j \neq i$ , received from the other subscribers according to  
 $k := f(z_1, \dots, z_n)$ ,  
 $f$  being a symmetrical function which is invariant under a permutation of its arguments.

4. (Withdrawn) The method as recited in claim 3 wherein the transmission key  $k^{ij}$  is known to subscriber  $T_j$  according to  $k^{ij} = k^{ji}$ .

5. (Currently Amended) A method for establishing a common key for a group of at least three subscribers for transmitting messages over a communication channel, the method comprising the steps of:

generating, by each subscriber  $T_j$ , a respective message  $N_j = (g^{z_j} \bmod p)$  from a publicly known element  $g$  of large order of a publicly known mathematical group  $G$  and a respective random number  $z_j$ ,  $j = 1$  to  $n$ , where  $n$  is the number of subscribers in the group of at least three subscribers;

sending the respective message  $N_j$  by each subscriber ~~except a predetermined first subscriber  $T_1$  of the at least three subscribers,~~ to each of the other subscribers  $T_j$ ,  $j \neq i$  the first subscriber  $T_i$ ;

encrypting, by the a first subscriber  $T_1$ , the received messages  $N_j$  of the other subscribers  $T_j$ ,  $j \neq 1$ , with the random number  $z_1$  to form a respective transmission key  $k^{1j}$  for each subscriber  $T_j$ ,  $j \neq 1$ ;

computing, by each subscriber  $T_j$ ,  $j \neq 1$ , a symmetrical counterpart  $k^{j1}$  of the respective transmission key  $k^{1j}$  using the received message  $N_1$ ;

sending, by the first subscriber  $T_1$ , the random number  $z_1$  to all other subscribers  $T_j$ ,  $j \neq 1$  in encrypted form by generating a message  $M_{1j}$  according to  $M_{1j} := E(k^{1j}, z_1)$ ,  $E(k^{1j}, z_1)$  being a symmetrical encryption algorithm in which the random number  $z_1$  is encrypted with the transmission key  $k^{1j}$ ; and

decrypting, by each subscriber  $T_j$ , the message  $M_{1j}$ ;

determining a common key  $k$ , by each subscriber  $T_j$ , using an assignment  $k := h(z_1, g^{z_2}, \dots, g^{z_n})$ ,  $h(x_1, x_2, \dots, x_n)$  being a function which is symmetrical in the arguments  $x_2, \dots, x_n$ ;

encrypting, by one of the subscribers  $T_i$ , a transmission message using the common key  $k_i$  and being useable for transmitting messages over a communication channel  
transmitting the encrypted transmission message to at least one of the other subscribers  $T_j, j \neq i$ .

6. (Previously Presented) The method as recited in claim 5 wherein the transmission key is known to subscriber  $T_j$  according to  $k^{ij} = k^{ji}$ .

7. (Currently Amended) A method for establishing a common key for a group of subscribers for encryption and decryption of messages, the method comprising the steps of:

each of the subscribers  $T_j$  generating a respective random number  $z_j$ , where  $j$  goes from 1 to  $n$  and  $n$  is the number of subscribers in the group of subscribers;

each of the subscribers  $T_j$  generating a respective first message  $N_j = (g^{z_j} \bmod p)$  from a publicly known element  $g$  of large order of a publicly known mathematical group  $G$ ;

each of the subscribers  $T_j, j \neq 1$ , sending the respective first message  $N_j$  to a first subscriber  $T_1$ ; each of the other subscribers  $T_j, j \neq i$ ;

the a first subscriber  $T_1$  computing a transmission key  $k^{ij} = N_j^{z_1} \bmod p$  for each of the other subscribers  $T_j, j \neq 1$ , based on the received respective first message  $N_j, j \neq 1$ ;

each of the subscribers  $T_j, j \neq 1$ , computing a symmetrical counterpart  $k^{ji}$  of the respective transmission key  $k^{ij}$  using the received first message  $N_i$ ;

the first subscriber  $T_1$  encrypting a second message  $M_{1j} := E(k^{1j}, z_1)$  for each of the other subscribers  $T_j$ ,  $j \neq 1$ , where  $E(k^{1j}, z_1)$  is a symmetrical encryption algorithm in which  $z_1$  is encrypted with the transmission key  $k^{1j}$ ;

the first subscriber  $T_1$  sending the encrypted second message  $M_{1j}$  to each of the other subscribers  $T_j$ ,  $j \neq 1$ ; and

each of the subscribers  $T_j$  decrypting the second message  $M_{1j}$ ;

each of the subscribers  $T_i$  computing a common key  $k$  according to an assignment  $k := h(z_1, g^{z_2}, \dots, g^{z_n})$ , where  $h(x_1, x_2, \dots, x_n)$  is a symmetrical function;

a subscriber  $T_i$  encrypting a third message using the common key  $k$ ; and

the subscriber  $T_i$  transmitting the encrypted third message to at least one of the other subscribers  $T_j$ ,  $j \neq i$ ;

8. (Currently Amended) The method according to claim 7, wherein the each respective random number  $z_j$  is selected from the set  $\{1, \dots, p-2\}$ .

9. (Previously Presented) The method according to claim 7, wherein the length of  $p$  is at least 1024 bits.

10. (Previously Presented) The method according to claim 7, wherein  $g$  has a multiplicative order of at least  $2^{160}$ .

11. (Previously Presented) The method according to claim 7 wherein the transmission key is known to a respective subscriber  $T_j$  according to  $k^{ij} = k^{ji}$ .

12. (Previously Presented) The method according to claim 7, wherein  $h(z1, g^{z2}, \dots, g^{zn}) = g^{z1 \cdot z1} * g^{z2 \cdot z1} * \dots * g^{zn \cdot z1}$ .

13. (Currently Amended) A method for establishing a common key for a group of subscribers for encryption and decryption of messages, the method comprising the steps of:

each of the subscribers  $T_j$  generating a respective random number  $z_j$ , where  $j$  goes from 1 to  $n$  and  $n$  is the number of subscribers in the group of subscribers;

each of the subscribers  $T_j$  storing the respective random number  $z_j$  in a respective memory;

each of the subscribers  $T_j$  generating a respective first message  $N_j = (g^{z_j} \bmod p)$  from a publicly known element  $g$  of large order of a publicly known mathematical group  $G$ ;

each of the subscribers  $T_j, j \neq 1$ , sending the respective first message  $N_j, j \neq 1$  to a first subscriber  $T_1$ , each of the other subscribers  $T_j, j \neq 1$ ;

the first subscriber  $T_1$  storing each of the received first messages  $N_j, j \neq 1$  in a memory;

the first subscriber  $T_1$  computing a transmission key  $k^{ij} = N_j^{z1} \bmod p$  for each of the other subscribers  $T_j, j \neq 1$ , based on the received respective first message  $N_j, j \neq 1$ ;

each of the subscribers  $T_j, j \neq 1$ , computing a symmetrical counterpart  $k^{ji}$  of the respective transmission key  $k^{ij}$  using the received first message  $N_j$ ;



~~the one of the respective subscribers~~ subscriber  $T_i$  transmitting the encrypted third message to each of the other respective subscribers  $T_j, j \neq i$ ;

each of the other respective subscribers  $T_j, j \neq i$  decrypting the received encrypted third message using the computed common  $k$ .

16. (New) The method according to claim 5, further comprising the step of:

decrypting, by the at least one of the other subscribers  $T_j, j \neq i$ , the transmitted transmission message using the common key  $k$ .

17. (New) The method according to claim 7, further comprising the step of:

the at least one other subscriber  $T_j, j \neq i$  decrypting the received third message using the common key  $k$ .

18. (New) The method according to claim 13, further comprising the step of:

the at least one other subscriber  $T_j, j \neq i$  decrypting the received third message using the common key  $k$ .